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Note

The 3-choosability of plane graphs of girth 4<sup>☆</sup>Peter C.B. Lam<sup>a</sup>, Wai Chee Shiu<sup>a</sup>, Zeng Min Song<sup>b</sup><sup>a</sup>Department of Mathematics, Hong Kong Baptist University, Hong Kong, China<sup>b</sup>Department of Applied Mathematics, Southeast University, Nanjing, China

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**Abstract**

A set  $S$  of vertices of the graph  $G$  is called  $k$ -reducible if the following is true:  $G$  is  $k$ -choosable if and only if  $G - S$  is  $k$ -choosable. A  $k$ -reduced subgraph  $H$  of  $G$  is a subgraph of  $G$  such that  $H$  contains no  $k$ -reducible set of some specific forms. In this paper, we show that a 3-reduced subgraph of a non-3-choosable plane graph  $G$  contains either adjacent 5-faces, or an adjacent 4-face and  $k$ -face, where  $k \leq 6$ . Using this result, we obtain some sufficient conditions for a plane graph to be 3-choosable. In particular, if  $G$  is of girth 4 and contains no 5- and 6-cycles, then  $G$  is 3-choosable. © 2005 Published by Elsevier B.V.

MSC: 05C15; 05C78

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**1. Introduction**

In this paper, we consider only finite and simple graphs. Undefined terms may be found in [2]. Suppose  $k$  is an integer. Then  $k^+$  and  $k^-$  denote integers  $\geq k$  and  $\leq k$ , respectively. A vertex  $u$  is called a  $k$ -vertex if  $d_G(u) = k$ . If  $k = 3$ , then  $u$  is called a *minor* vertex. A face  $f$  is called a  $k$ -face if  $d_G(f) = k$ . If  $k = 4$ , then  $f$  is called a *minor* face. If no confusion can arise,  $d(v)$  and  $d(f)$  will be used instead of  $d_G(v)$  and  $d_G(f)$ , respectively. A face of a plane graph is *incident* with all edges and vertices on its boundary. Two faces are *adjacent*

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if they have an edge in common. Let  $F_j$  be the set of  $j$ -faces incident with  $j$  minor vertices and  $F_j^i$  be the subset of  $F_j$  consisting only of faces adjacent to  $i$  minor faces.

A  $k$ -cycle is a cycle of length  $k$ . The name of a cycle will also be used to denote the set of its vertices. We use  $F(u)$  and  $k(u)$  to denote the set of all faces incident with the vertex  $u$  and  $\min\{d(x) | x \in F(u)\}$ , respectively.

A *list coloring* of  $G$  is an assignment of colors to  $V(G)$  such that each vertex  $v$  receives a color from a prescribed list  $L(v)$  of colors and adjacent vertices receive distinct colors (see [7]).  $L(G) = (L(v) | v \in V(G))$  is called a *color list* of  $G$ . The graph  $G$  is called  *$k$ -choosable* if  $G$  admits a list coloring for all color lists  $L$  with  $k$  colors in each list.

All 2-choosable graphs have been characterized by Erdős et al. [3]. In [5], Thomassen proved that every plane graph is 5-choosable. From that proof, one can find a simple linear algorithm for finding the corresponding list coloring. Voigt [8] showed that there are planar graphs which are not 4-choosable. It remains to decide whether a given plane graph is 4- or 3-choosable. Gutner [4] proved that these problems are NP-hard. So far, some sufficient conditions have been obtained and some constructions have been found. Alon and Tarsi [1] proved that every plane bipartite graph is 3-choosable. Thomassen [6] proved that every plane graph of girth at least 5 is 3-choosable. Voigt [9] gave an example of a non-3-choosable plane graph of girth 4. The objective of this paper is to study 3-choosability of plane graphs of girth 4. In this paper we give a necessary condition for a plane graph of girth 4 to be 3-choosable. In particular, if  $G$  is of girth 4 and contains no 5- and 6-cycles, then  $G$  is 3-choosable.

## 2. Some lemmas

Let  $C$  be a cycle in  $G$ . A *sub-cycle* of  $C$  is a cycle in  $G[C]$  with length strictly shorter than  $C$ . We shall need some lemmas, the first of which is well known.

**Lemma 1.** *Let  $v$  be a vertex of  $G$ , not necessarily planar, with  $d(v) \leq k - 1$ . Then  $G$  is  $k$ -choosable if and only if  $G - \{v\}$  is  $k$ -choosable.*

**Lemma 2.** *Let  $G$  be a graph having no subgraphs isomorphic to  $K_4$ . Suppose  $n \geq 2$  and  $C$  is a  $2n$ -cycle of  $G$  and  $d(u) \leq k$  for all  $u \in C$ . Then either (i) there exists an even sub-cycle  $C^*$  of  $C$  and with  $d(u) \leq k$  for all  $u \in C^*$  or (ii)  $G$  is  $k$ -choosable if and only if  $G - C$  is  $k$ -choosable.*

**Proof.** Suppose that  $G - C$  is  $k$ -choosable and suppose that  $L = (L(v) | v \in V(G))$  is a color list of  $G$  in which each list contains  $k$  colors. Let  $\phi$  be an  $L$ -list coloring of  $G - C$ . For all  $v \in C$ , let  $L^0(v) = L(v) \setminus \{\phi(u) | u \in V(G) \setminus C \text{ and } vu \in E(G)\}$ . Then  $|L^0(v)| \geq 2$ . We have three cases:

*Case 1:*  $G[C]$  is chordless. Since even cycles are 2-choosable, there exists a 2-list coloring  $\phi'$  of  $C$ . Combining  $\phi$  and  $\phi'$ , we obtain an  $L$ -list coloring of  $G$ .

*Case 2:*  $G[C]$  contains exactly one chord. Without loss of generality, assume that  $C = u_1 u_2 \cdots u_{2n}$  and  $u_1 u_j$  is the chord, where  $3 \leq j \leq 2n - 1$ . Then  $|L^0(u_1)| \geq 3$ ,  $|L^0(u_j)| \geq 3$  and  $|L^0(u_i)| \geq 2$  if  $i \neq 1$  or  $j$ . We can then define  $\phi'$  on  $C$  as follows: choose  $\phi'(u_j) \in L^0(u_j) \setminus L^*$ , where  $L^*$  is a subset of  $L^0(u_{j-1})$  with  $|L^*| = 2$ ,  $\phi'(u_i) \in L^0(u_i) \setminus \{\phi'(u_{i-1})\}$  for

$i = j + 1, j + 2, \dots, 2n$ ,  $\phi'(u_1) \in L^0(u_1) \setminus \{\phi'(u_{2n}), \phi'(u_j)\}$ ,  $\phi'(u_i) \in L^0(u_i) \setminus \{\phi'(u_{i-1})\}$  for  $i = 2, \dots, j - 2$  (if  $j \geq 4$ ) and  $\phi'(u_{j-1}) \in L^* \setminus \{\phi'(u_{j-2})\}$ . Combining  $\phi$  and  $\phi'$ , we obtain an  $L$ -list coloring of  $G$ .

*Case 3:*  $G[C]$  has two or more chords. If  $n = 2$ , then  $G[C]$  is isomorphic to  $K_4$ , contrary to the conditions of the lemma. So we may assume that  $n \geq 3$ . Suppose  $u_1u_i$  and  $u_ju_k$  are the two chords. If  $3 \leq i \leq j < k - 1 \leq 2n - 1$ , then  $u_1u_2 \cdots u_i, u_1u_iu_{i+1} \cdots u_ju_ku_{k+1} \cdots u_{2n}$  and  $u_ku_ju_{j+1} \cdots u_{k-1}$  are also three cycles with total length of  $2n + 4$ . Therefore, one of these cycles must be even. Since the length of each cycle must be at least 3, length of each of these cycles is at most  $2n - 2$ . If  $2 \leq j < i < k \leq 2n$ , then the length of the cycles  $C_1 = u_1u_2 \cdots u_i$ ,  $C_2 = u_ju_{j+1} \cdots u_k$  and  $C_3 = u_1 \cdots u_ju_ku_{k-1} \cdots u_i$  are  $i, k - j + 1$  and  $j + k - i + 1$ , respectively. The total length of  $C_1, C_2$  and  $C_3$  is therefore even. If either  $C_1$  or  $C_2$  is even, then we have obtained an even sub-cycle of  $C$ . If both  $C_1$  and  $C_2$  are odd, then  $C_3$  is even. If  $j + k - i + 1 = 2n$ , then  $C_4 = u_1u_iu_{i-1} \cdots u_ju_ku_{k+1} \cdots u_{2n}$  is of length  $2n - k + j - i + 3 = 4$ . Therefore either  $C_3$  or  $C_4$  is an even sub-cycle of  $C$ .  $\square$

A set  $S$  of vertices of graph  $G$  is called  $k$ -reducible if the following is true:  $G$  is  $k$ -choosable if and only if  $G - S$  is  $k$ -choosable. Hence by Lemmas 1 and 2, if  $\delta(G) \leq k - 1$ , or if  $G$  has an even cycle passing through  $k^-$ -vertices, then  $G$  contains a  $k$ -reducible subset. A  $k$ -reduced subgraph of  $G$  is obtained from  $G$  by deleting successively  $k$ -reducible sets as determined by Lemmas 1 and 2 until no such set exists. If  $G$  contains no  $k$ -reducible sets as determined by Lemmas 1 and 2, we say that  $G$  is  $k$ -reduced. The following lemma follows from definition.

**Lemma 3.**  $G$  is  $k$ -choosable if and only if a  $k$ -reduced subgraph of  $G$  is  $k$ -choosable.

By Lemma 3, in order to determine whether a graph  $G$  is 3-choosable or not, it is sufficient to consider the choosability of a 3-reduced subgraph of  $G$ . The following lemma can be obtained from definition and Lemma 2.

**Lemma 4.** If  $H$  is a  $k$ -reduced subgraph of  $G$  and  $H$  does not contain  $K_4$ 's, then any even cycle of  $H$  contains a  $(k + 1)^+$ -vertex.

### 3. Main results

**Theorem 5.** Let  $G$  be a non-3-choosable plane graph of girth not less than 4 and let  $H$  be any 3-reduced subgraph of  $G$ . Then there exists in  $H$  adjacent 5-faces or adjacent 4-face and  $k$ -face ( $k \leq 6$ ).

**Corollary 6.** Let  $G$  be a plane graph of girth not less than 4. If a 3-reduced subgraph of  $G$  contains neither adjacent 5-faces nor adjacent 4-face and  $k$ -face ( $k \leq 6$ ), then  $G$  is 3-choosable.

**Corollary 7.** Let  $G$  be a plane graph of girth not less than 4. If  $G$  contains neither adjacent 5-cycles nor adjacent 4-cycle and  $k$ -cycle ( $k \leq 6$ ), then  $G$  is 3-choosable.

Clearly, graphs satisfying conditions of Corollary 7 may contain 4-, 5- or 6-cycles.

**Corollary 8.** *Let  $G$  be a plane graph of girth not less than 4. If  $G$  contains no 5- and 6-cycles, then  $G$  is 3-choosable.*

**Proof.** Let  $G$  be a plane graph without 5- and 6-cycles. Consider a 3-reduced subgraph  $H$  of  $G$ . Certainly  $\delta(H) \geq 3$  and  $H$  contains no 5- and 6-cycles. Hence,  $H$  contains neither adjacent 5-faces nor adjacent 4-face and  $k$ -face ( $k \leq 6$ ). It follows from Corollary 6 that  $H$  is 3-choosable. Thus, the corollary holds by Lemma 3.  $\square$

#### 4. Proof of Theorem 5

Let  $G$  be a non-3-choosable plane graph. Without loss of generality, we assume that  $G$  is 3-reduced. Suppose the theorem is false. Then there exist neither adjacent 5-faces nor adjacent 4-face and  $k$ -face with  $k \leq 6$ . That is, there are no adjacent 5-faces and  $d(f) \geq 7$  for every face  $f$  adjacent to a 4-face. This also implies that  $F_7 = F_7^0 \cup F_7^1 \cup F_7^2 \cup F_7^3$ . If we assign a weight of  $\sigma(x) = d(x) - 4$  to each  $x \in V(H) \cup F(H)$ , then by Euler's formula we have:

$$\sum_{v \in V(H) \cup F(H)} \sigma(x) = -8. \quad (1)$$

If we obtain a new weight  $\sigma^*(x)$  for each  $x \in V(H) \cup F(H)$  by transferring weights from one element to another, then we also have

$$\sum_{v \in V(H) \cup F(H)} \sigma^*(x) = -8. \quad (2)$$

Moreover, if  $\sigma^*(x) \geq 0$  for all  $x \in V(H) \cup F(H)$ , then the theorem is proved.

Weights will be transferred from  $f \in F$  to  $v \in V$ , where  $v$  is incident with  $f$ ,  $d(f) \geq 5$  and  $d(v) = 3$ , according to the following rules:

- (R1) Transfer  $\frac{1}{2}$  if (a)  $d(f) \geq 7$  and  $f \notin F_7$  or (b)  $f \in F_7$  and  $k(v) = 4$ .
- (R2) Transfer  $\frac{2}{5}$  if (a)  $d(f) = 6$  or (b)  $f \in F_7$  and  $k(v) = 5$ .
- (R3) Transfer  $\frac{1}{5}$  if (a)  $d(f) = 5$  or (b)  $f \in F_7$  and  $k(v) = 6$ .

Let  $v \in V$ . If  $d(v) \geq 4$ , then  $\sigma^*(v) = \sigma(v) \geq 0$ . Suppose  $d(v) = 3$ ,  $v$  is incident with  $f_1, f_2$  and  $f_3$  and  $d(f_1) \leq d(f_2) \leq d(f_3)$ . By Lemma 4, at most one of these three faces is in  $F_7$ . If  $d(f_1) = 4$ , then (R1) is applicable to both  $f_2$  and  $f_3$ . Therefore,  $\sigma^*(v) = \sigma(v) + 2 \cdot \frac{1}{2} = 0$ . If  $d(f_1) = 5$ , then (R3(a)) is applicable to  $f_1$  and (R3) is not applicable to  $f_2$  and  $f_3$ . So  $\sigma^*(v) \geq \sigma(v) + \frac{1}{5} + 2 \cdot \frac{2}{5} = 0$ . If  $d(f_1) = 6$ , then (R3(b)) is applicable to at most one of  $f_2$  and  $f_3$ , and (R3(a)) is not applicable. So  $\sigma^*(v) \geq \sigma(v) + \frac{1}{5} + 2 \cdot \frac{2}{5} = 0$ . If  $d(f_1) \geq 7$ , then (R1(a)) is applicable to at least two of  $f_1, f_2$  and  $f_3$ . So  $\sigma^*(v) \geq \sigma(v) + 2 \cdot \frac{1}{2} = 0$ .

Let  $f \in F$ . If  $d(f) = 4$ , then  $\sigma^*(f) = \sigma(f) = 0$ . If  $d(f) = 5$ , then  $\sigma^*(f) \geq \sigma(f) - 5 \cdot \frac{1}{5} = 0$ . If  $d(f) = 6$ , then at least one  $4^+$ -vertex is incident with  $f$  by Lemma 4. So  $\sigma^*(f) \geq$

$\sigma(f) - 5 \cdot \frac{2}{5} = 0$ . If  $d(f) = k \geq 8$ , then  $\sigma^*(f) \geq \sigma(f) - k \cdot \frac{1}{2} \geq k - 4 - (k/2) \geq 0$ . If  $d(f) = 7$  and  $f \notin F_7$ , then at least one  $4^+$ -vertex is incident with  $f$ . So  $\sigma^*(f) \geq \sigma(f) - 6 \cdot \frac{1}{2} = 0$ .

If  $f \in F_7^0 \cup F_7^1$ , then weights are transferred from  $f$  to at most two 3-vertices according to [R1(b)]. So  $\sigma^*(f) \geq \sigma(f) - 2 \cdot \frac{1}{2} - 5 \cdot \frac{2}{5} = 0$ .

If  $f \in F_7^2$  and  $f$  is not adjacent to any 6-face, then at least one vertex incident with  $f$ , say  $v$ , is incident with  $7^+$ -faces only, and 0 is transferred from  $f$  to  $v$ . So  $\sigma^*(f) \geq \sigma(f) - 4 \cdot \frac{1}{2} - 2 \cdot \frac{2}{5} > 0$ . If  $f \in F_7^2$  and  $f$  is adjacent to a 6-face  $f_1$ , then at least one face adjacent to both  $f$  and  $f_1$  is a  $7^+$ -face. So  $\sigma^*(f) \geq \sigma(f) - 4 \cdot \frac{1}{2} - 2 \cdot \frac{2}{5} - \frac{1}{5} = 0$ .

If  $f \in F_7^3$ , then one vertex incident with  $f$  is incident with  $7^+$ -faces only, and the weight transferred from  $f$  to  $v$  is 0. So  $\sigma^*(f) \geq \sigma(f) - 6 \cdot \frac{1}{2} = 0$ .

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